Multicut and integral multiflow : a survey.

Marie-Christine Costa, Lucas Létocart, Frédéric Roupin

Erratum to "M.-C. Costa, L. Létocart and F. Roupin. Minimal multicut and maximal integer multiflow: a survey" [European Journal of Operational Research 162(1) (2005), 55-69]

C. Bentz [a], M.-C. Costa [a,*], L. Létocart [b], F. Roupin [c]

[a] CEDRIC, CNAM, 292 rue Saint Martin, 75141 Paris cedex 03, France
[b] LIPN, Institut Galilée, Université Paris XIII, avenue J.B. Clément, 93430 Villetaneuse, France
[c] CEDRIC, CNAM-IIE, 18 allée Jean Rostand, 91025 Evry cedex, France

Page 56, first paragraph (also page 59, beginning of the second paragraph on the right column; Page 60, beginning of the last paragraph on the right column; Page 67, first paragraph on the right column): actually, for these problems, the general directed case is "harder" than the undirected one, because there exists a linear reduction from the latter to the former (see (70.9) on p. 1224 in [A. Schrijver. Combinatorial Optimization - Polyhedra and Efficiency. Algorithms and Combinatorics 24 (2003). Springer]).

Page 59, on the top of the right column: the algorithm in [26] does not provide an approximate solution for IMFP.

Page 61, first paragraph on the left column (also page 66 (table), line "Directed graphs", column "Multiterminal cut"): actually, the algorithm of Naor and Zosin [50] provides a 2-approximation for the MULTITERMINAL CUT problem in directed graphs.

Page 61, Section 3.4.2 (also page 62, last paragraph before Section 4.2; Page 66 (table), line "Undirected graphs", column "EdgeDisjPath"; Page 67, end of the first paragraph on the right column): actually, although the *edge-disjoint paths problem* is tractable for fixed K in undirected graphs [55] and in acyclic digraphs [22], the MAX EDGEDISJPATH problem is NP-Hard in undirected graphs and in acyclic digraphs for $K = 2$ [21].

Page 65, first paragraph of Section 4.4: actually, in [59], Sebö proved that the demand version of IMFP is polynomial-time solvable in this case, but not IMFP.

Page 66 (table), line "Undirected graphs", column "IMFP": actually, it is known that IMFP is NP-Hard to approximate within $m^{\frac{1}{2}-\epsilon}$ in directed graphs only, while in undirected graphs a recent (and weaker) result is that it is NP-Hard to approximate within $(\log m)^{\frac{1}{3}-\epsilon}$ [M. Andrews and L. Zhang. Hardness of the undirected edge-disjoint paths problem. Proceedings STOC'05 (2005)].

*Corresponding author. Tel.: +33-140-272524; Fax: +33-140-272296
E-mail adresses: cedric.bentz@cnam.fr (C. Bentz), costa@cnam.fr (M.-C. Costa), lucas.letocart@lipn.univ-paris13.fr (L. Létocart), roupin@iie.cnam.fr (F. Roupin).